Formulation 000	Numerical scheme	Complexity 00	Memory reductio

A probabilistic numerical method for optimal multiple switching problem and application to investments in electricity generation

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Motivation							

#### Investments in power generation

- Capital intensive
- Very long-term returns
- Many technologies
- Many random factors :
  - Demand
  - Outages
  - Fuel prices
  - Inflows

#### Natural modelling : Stochastic control

- High-dimensional
- Infinite horizon

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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Motivat	ion				

More precisely : Optimal multiple switching

#### Outline

- 1 Numerical scheme
  - Convergence rate
  - Complexity analysis

# 2 Application

- Detailed modelling
- Numerical solution

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Outline					

#### 1 Formulation

2 Numerical scheme

# 3 Complexity

- 4 Memory reduction
- **5** Application

### 6 Conclusion

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Formulation ●00	Numerical scheme	Complexity 00	Memory reduction	Application 0000000000000	Conclusion 0
Problem					

$$v(t,x,i) = \sup_{\alpha \in \mathcal{A}_{t,i}} \mathbb{E}\left[\int_{t}^{\infty} f(s, X_{s}^{t,x}, I_{s}^{\alpha}) ds - \sum_{\tau_{n} \geq t} k(\tau_{n}, \zeta_{n})\right]$$

•  $X^{t,x}$  Markovian diffusion in  $\mathbb{R}^d$ ,  $X_t = x$ 

- $I^{\alpha}$  piecewise constant in a finite set  $\mathbb{I}_q = \{i_1, i_2, \dots, i_q\}$ ( $\Rightarrow$  optimal switching)
- α = (τ<sub>n</sub>, ι<sub>n</sub>)<sub>n∈N</sub> impulse control strategy, τ<sub>n</sub> increasing stopping times, ι<sub>n</sub> random variables in I<sub>q</sub>, ζ<sub>n</sub> = ι<sub>n</sub> − ι<sub>n-1</sub>

$$I_s^{\alpha} = \iota_0 \mathbf{1} \{ 0 \le s < \tau_0 \} + \sum_{n \in \mathbb{N}} \iota_n \mathbf{1} \{ \tau_n \le s < \tau_{n+1} \}$$

•  $\mathcal{A}_{t,i}$  set of admissible strategies s.t.  $I_t^{\alpha} = i$ 

Formulation 0●0	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Assump	tions				

#### Diffusion coefficients

Lipschitz continuity & Linear growth

#### Reward function f

Lipschitz continuity & Linear growth Exponential discount  $e^{-\rho t}$ 

#### Cost function k

Lipschitz continuity & Linear growth Exponential discount  $e^{-\rho t}$ Minimum fixed cost Triangular condition : k(t; i, j) + k(t; j, k) > k(t; i, k)

Formulation	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Properti	es				

- v<sub>i</sub> = v (.,.,i) solutions of a system of Hamilton-Jacobi-Bellman Quasi-Variational Inequalities
- v(t, X<sub>t</sub>, i) solution of a Reflected Backward Stochastic Differential Equation
- **Dynamic Programming Principle** :  $\forall \tau \geq t$  stopping time,

$$v(t,x,i) = \sup_{\alpha \in \mathcal{A}_{t,i}} \mathbb{E}\left[\int_{t}^{\tau} f(s, X_{s}^{t,x}, I_{s}^{\alpha}) ds - \sum_{t \leq \tau_{n} \leq \tau} k(\tau_{n}, \zeta_{n}) + v(\tau, X_{\tau}^{t,x}, I_{\tau}^{\alpha})\right]$$

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Outline					

#### **1** Formulation

#### 2 Numerical scheme

# 3 Complexity

4 Memory reduction

# 5 Application

#### 6 Conclusion

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Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
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Finite t	ime horizon				

$$v_{T}(t,x,i) = \sup_{\alpha \in \mathcal{A}_{t,i}^{T}} \mathbb{E}\left[\int_{t}^{T} f(s,X_{s}^{t,x},I_{s}^{\alpha}) ds - \sum_{t \leq \tau_{n} \leq T} k(\tau_{n},\zeta_{n}) + g(T,X_{T}^{t,x},I_{T}^{\alpha})\right]$$

$$g(T, x, i) = \mathbb{E}\left[\int_{T}^{\infty} f\left(s, X_{s}^{T, x}, i\right) ds\right]$$
  
OR any Lipschitz fct s.t.  $|g(T, x, i)| \leq Ce^{-\rho T} (1 + |x|)$ 

$$|v(t,x,i) - v_T(t,x,i)| \le \mathsf{C} \left(1 + |\mathsf{x}|\right) \mathrm{e}^{-\bar{
ho}\mathsf{T}}$$

Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
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Time d	iscretization				

$$\bar{\nu}_{\Pi}(t,x,i) = \sup_{\alpha \in \mathcal{A}_{t,i}^{\Pi}} \mathbb{E}\left[\int_{t}^{T} f(\pi(s),\bar{X}_{s}^{t,x},I_{s}^{\alpha}) ds - \sum_{t \leq \tau_{n} \leq T} k(\tau_{n},\iota_{n-1},\iota_{n}) + g(T,\bar{X}_{T}^{t,x},I_{T}^{\alpha})\right]$$

$$\begin{aligned} \Pi &= \{ t_0 = 0 < t_1 < \ldots < t_N = T \}, \text{ step } h \\ d\bar{X}_s &= b\big(\pi(s), \bar{X}_{\pi(s)}\big) ds + \sigma\big(\pi(s), \bar{X}_{\pi(s)}\big) dW_s \,, \, 0 \leq s \leq T, \, \bar{X}_0 = x_0 \end{aligned}$$

$$|v_{\mathcal{T}}(0,x,i)-ar{v}_{\mathsf{\Pi}}(0,x,i)| \leq \mathsf{C}\left(1+|\mathbf{x}|^{rac{3}{2}}
ight)\sqrt{\mathsf{h}}$$

cf. [Gassiat et al., 2011]

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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Space I	ocalization				

$$\begin{split} \bar{X}_t \in \mathbb{R}^d &\to \mathcal{P}^{\varepsilon} \left( \bar{X}_t \right) \in \mathcal{D}^{\varepsilon} \\ \mathcal{D}^{\varepsilon} \text{ s.t. } \mathbb{E} \left[ \left| \bar{X}_t - \mathcal{P}^{\varepsilon} \left( \bar{X}_t \right) \right| \right] \leq \varepsilon \ \forall 0 \leq t \leq T \\ \mathcal{D}^{\varepsilon} \text{ bounded domain } : \forall x \in \mathcal{D}^{\varepsilon}, \ |x| \leq C \left( T, \varepsilon \right) \end{split}$$

$$\left| ar{v}_{\Pi} \left( 0,x,i 
ight) - ar{v}_{\Pi}^{arepsilon} \left( 0,x,i 
ight) 
ight| \leq \mathbf{C} arepsilon$$

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Dynamic programming principle

$$\begin{split} \bar{v}_{\Pi}\left(T, x, i\right) &= g\left(T, x, i\right) \\ \bar{v}_{\Pi}\left(t_{n}, x, i\right) &= \max_{j \in \mathbb{I}_{q}} \left\{ hf\left(t_{n}, x, j\right) - k\left(t_{n}, i, j\right) + \Phi_{j}^{t_{n}, x}\left(\bar{v}_{\Pi}\right) \right\} \\ \Phi_{j}^{t_{n}, x}\left(\varphi\right) &= \mathbb{E}\left[\varphi\left(t_{n+1}, \bar{X}_{t_{n+1}}, j\right) \left| \bar{X}_{t_{n}} = x \right] \end{split}$$

pproximation : Least-squares regression

$$\Phi_{j}^{t_{n},x}(\varphi) \rightsquigarrow \tilde{\Phi}_{j}^{t_{n},x}(\varphi) = \min_{\lambda \in \mathbb{R}^{K}} \mathbb{E}\left[ \left( \varphi\left(t_{n+1}, \bar{X}_{t_{n+1}}, i\right) - \sum_{k=1}^{K} \lambda_{k} e_{k}\left(\bar{X}_{t_{n}}\right) \right)^{2} \right]$$

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Dynamic programming principle

$$\begin{split} \bar{v}_{\Pi}\left(T, x, i\right) &= g\left(T, x, i\right) \\ \bar{v}_{\Pi}\left(t_{n}, x, i\right) &= \max_{j \in \mathbb{I}_{q}} \left\{ hf\left(t_{n}, x, j\right) - k\left(t_{n}, i, j\right) + \Phi_{j}^{t_{n}, x}\left(\bar{v}_{\Pi}\right) \right\} \\ \Phi_{j}^{t_{n}, x}\left(\varphi\right) &= \mathbb{E}\left[\varphi\left(t_{n+1}, \bar{X}_{t_{n+1}}, j\right) \left| \bar{X}_{t_{n}} = x \right] \end{split}$$

Approximation : Least-squares regression

$$\Phi_{j}^{t_{n},x}(\varphi) \rightsquigarrow \tilde{\Phi}_{j}^{t_{n},x}(\varphi) = \min_{\lambda \in \mathbb{R}^{K}} \mathbb{E}\left[ \left( \varphi\left(t_{n+1}, \bar{X}_{t_{n+1}}, i\right) - \sum_{k=1}^{K} \lambda_{k} e_{k}\left(\bar{X}_{t_{n}}\right) \right)^{2} \right]$$

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Local basis

$$e_k(x) = \mathbf{1}\{x \in B_k\}$$

 $(B_k)_{1 \leq k \leq K}$  hypercubes  $\subset \mathbb{R}^d$ , partition of  $\mathcal{D}^{\varepsilon}$ , edges of length  $\delta$ cf. [Gobet et al., 2005], [Bouchard and Warin, 2011]

Approximation error

$$\left| ar{v}_{\mathsf{\Pi}} \left( 0, x, i 
ight) - ar{v}_{\mathsf{\Pi}} \left( 0, x, i 
ight) 
ight| \leq \mathbf{C} rac{\delta}{\mathbf{h}}$$

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Conditio	nal expecta	tion app	roximation (	(2/2)	

#### Sample approximation

$$\tilde{\Phi}_{j}^{t_{n},x}(\varphi) \rightsquigarrow \hat{\Phi}_{j}^{t_{n},x}(\varphi) = \min_{\lambda \in \mathbb{R}^{K}} \sum_{m=1}^{M} \left[ \left( \varphi\left(t_{n+1}, \bar{X}_{t_{n+1}}^{m}, i\right) - \sum_{k=1}^{K} \lambda_{k} e_{k}\left(\bar{X}_{t_{n}}^{m}\right) \right)^{2} \right]$$

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$$\forall p \geq 1, \ \left\| \hat{\Phi}_{j}^{t_{n,x}}(\varphi) - \tilde{\Phi}_{j}^{t_{n,x}}(\varphi) \right\|_{p} \leq \frac{C_{p}}{\sqrt{M}} \frac{\Gamma^{t_{n,x}}(\varphi) + \bar{\varphi}^{t_{n}}}{\mathbb{P}\left( \bar{X}_{t_{n}} \in B(x) \right)^{1 - \frac{1}{p \sqrt{2}}}}$$

- Extension of [Tan, 2011]
- Makes use of  $\left\|\frac{1}{M}\sum_{m=1}^{M} X_m\right\|_p \le \frac{C_p}{\sqrt{M}} \|X\|_{p\vee 2}$ (Marcinkiewicz-Zygmund + Jensen)

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Conditio	nal expecta	tion app	roximation (	(2/2)	

#### Sample approximation

$$\tilde{\Phi}_{j}^{t_{n},x}(\varphi) \rightsquigarrow \hat{\Phi}_{j}^{t_{n},x}(\varphi) = \min_{\lambda \in \mathbb{R}^{K}} \sum_{m=1}^{M} \left[ \left( \varphi\left(t_{n+1}, \bar{X}_{t_{n+1}}^{m}, i\right) - \sum_{k=1}^{K} \lambda_{k} e_{k}\left(\bar{X}_{t_{n}}^{m}\right) \right)^{2} \right]$$

#### Lemma

$$\forall p \geq 1 , \ \left\| \hat{\Phi}_{j}^{t_{n},x}\left(\varphi\right) - \tilde{\Phi}_{j}^{t_{n},x}\left(\varphi\right) \right\|_{p} \leq \frac{C_{p}}{\sqrt{M}} \frac{\Gamma^{t_{n},x}\left(\varphi\right) + \bar{\varphi}^{t_{n}}}{\mathbb{P}\left(\bar{X}_{t_{n}} \in B\left(x\right)\right)^{1 - \frac{1}{p \vee 2}}}$$

Extension of [Tan, 2011]
 Makes use of \$\|\frac{1}{M}\sum\_{m=1}^M X\_m \$\|\|\_p\$ ≤ \$\frac{C\_p}{\sqrt{M}}\$ \$\|X\|\_{p\vee 2}\$ (Marcinkiewicz-Zygmund + Jensen)

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Rate of	convergenc	е			

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$$\left\| \tilde{v}_{\Pi} \left( 0, x, i \right) - \hat{v}_{\Pi} \left( 0, x, i \right) \right\|_{p} \leq \frac{C_{p}}{h\sqrt{M}} \frac{1 + C(T, \varepsilon) + \sqrt{h}}{p(T, \delta, \varepsilon)^{1 - \frac{1}{p \vee 2}}}$$
$$P(T, \delta, \varepsilon) = \min_{t_{n} \in \Pi, B_{k} \subset \mathcal{D}^{\varepsilon}} \mathbb{P} \left( \bar{X}_{t_{n}} \in B_{k} \right)$$

#### Full rate of convergence

$$\|v(0,x,i) - \hat{v}_{\Pi}(0,x,i)\|_{p} \leq C_{p} \left\{ (1+|x|)e^{-\bar{\rho}T} + \left(1+|x|^{\frac{3}{2}}\right)\sqrt{h} + \varepsilon + \frac{\delta}{h} + \frac{C_{p}}{h\sqrt{M}} \frac{1+C(T,\varepsilon) + \sqrt{h}}{p(T,\delta,\varepsilon)^{1-\frac{1}{p\sqrt{2}}}} \right\}$$

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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion 0
Rate of	convergenc	е			

$$\left\| \tilde{v}_{\Pi} \left( 0, x, i \right) - \hat{v}_{\Pi} \left( 0, x, i \right) \right\|_{p} \leq \frac{C_{p}}{h\sqrt{M}} \frac{1 + C(T, \varepsilon) + \sqrt{h}}{p(T, \delta, \varepsilon)^{1 - \frac{1}{p \vee 2}}}$$
$$(T, \delta, \varepsilon) = \min_{t_{n} \in \Pi, B_{k} \subset \mathcal{D}^{\varepsilon}} \mathbb{P} \left( \bar{X}_{t_{n}} \in B_{k} \right)$$

# Full rate of convergence

$$\|v(0,x,i) - \hat{v}_{\Pi}(0,x,i)\|_{p} \leq C_{p} \left\{ (1+|x|)e^{-\bar{\rho}T} + \left(1+|x|^{\frac{3}{2}}\right)\sqrt{h} + \varepsilon + \frac{\delta}{h} + \frac{C_{p}}{h\sqrt{M}} \frac{1+C(T,\varepsilon) + \sqrt{h}}{p(T,\delta,\varepsilon)^{1-\frac{1}{p\vee2}}} \right\}$$

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Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
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Example	е				

 $X_t = W_t d$ -dim. B.M.

$$C(T,\varepsilon) < \sqrt{T\log\left(\frac{8T}{\pi\varepsilon^{\frac{2}{d}}}\right)}$$
$$p(T,\varepsilon,\delta) \gg \frac{\varepsilon}{(4T)^d} \delta^d$$

ne-parameter rate of convergence

$$\|v(0,x,i) - \hat{v}_{\Pi}(0,x,i)\|_{p} \leq C_{p}\left(1 + |x|^{\frac{3}{2}}\right)\sqrt{h}$$

$$T = \frac{1}{2\bar{\rho}} \ln\left(\frac{1}{\bar{h}}\right), \ \varepsilon = \sqrt{h}, \ \delta = h^{\frac{3}{2}}$$
$$M = \frac{1}{2\bar{\rho}} \left(\frac{2}{\bar{\rho}}\right)^{2d} \left(\ln\left(\frac{1}{\bar{h}}\right)\right)^{2d+1} \ln\left(\frac{4}{\pi\bar{\rho}} \frac{\ln\left(\frac{1}{\bar{h}}\right)}{h^{\frac{1}{d}}}\right) \frac{1}{h^{3(d+1)}}$$

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Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
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Example	е				

 $X_{t} = W_{t} \text{ d-dim. B.M.}$   $C(T, \varepsilon) < \sqrt{T \log\left(\frac{8T}{\pi \varepsilon^{\frac{2}{d}}}\right)}$   $p(T, \varepsilon, \delta) \gg \frac{\varepsilon}{(4T)^{d}} \delta^{d}$ 

One-parameter rate of convergence

$$\|v(0,x,i) - \hat{v}_{\Pi}(0,x,i)\|_{p} \leq C_{p}\left(1 + |x|^{\frac{3}{2}}\right)\sqrt{h}$$

$$T = \frac{1}{2\bar{\rho}} \ln\left(\frac{1}{\bar{h}}\right), \ \varepsilon = \sqrt{\bar{h}}, \ \delta = h^{\frac{3}{2}}$$
$$M = \frac{1}{2\bar{\rho}} \left(\frac{2}{\bar{\rho}}\right)^{2d} \left(\ln\left(\frac{1}{\bar{h}}\right)\right)^{2d+1} \ln\left(\frac{4}{\pi\bar{\rho}}\frac{\ln\left(\frac{1}{\bar{h}}\right)}{h^{\frac{1}{d}}}\right) \frac{1}{h^{3(d+1)}}$$

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Formulation	Numerical scheme	Complexity O	Memory reduction	Application 00000000000000	Conclusion O
Compu	tational com	plexity			

#### Dynamic programming

$$\begin{split} \hat{v}_{\Pi}\left(T, x_{m}, i\right) &= g\left(T, x_{m}, i\right) \\ \hat{v}_{\Pi}\left(t_{n}, x_{m}, i\right) &= \max_{j \in \mathbb{I}_{q}} \left\{ hf\left(t_{n}, x_{m}, j\right) - k\left(t_{n}, i, j\right) + \hat{\Phi}_{j}^{t_{n}, x_{m}}\left(\hat{v}_{\Pi}\right) \right\} \end{split}$$

# Complexity

$$\mathcal{O}\left(q^2 \times N \times M \log\left(M\right)\right)$$

- q =number of switches
- *N* =number of time steps
- *M* =number of Monte Carlo trajectories

or  $\mathcal{O}(q \times N \times M \log(M))$  under suitable conditions

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Memory	Complexity	/ : Naive	implementa	ation	

#### Euler scheme : Forward

$$\begin{aligned} x_{t_{n+1}} &= x_{t_n} + b(t_n, x_{t_n}) h + \sigma(t_n, x_{t_n}) \varepsilon_i \sqrt{h} \\ \varepsilon_i &\sim \mathcal{N}(0, 1) \end{aligned}$$

# Dynamic programming : Backward

$$\hat{v}_{\Pi}(T, x_{m}, i) = g(T, x_{m}, i)$$
$$\hat{v}_{\Pi}(t_{n}, x_{m}, i) = \max_{j \in \mathbb{I}_{q}} \left\{ hf(t_{n}, x_{m}, j) - k(t_{n}, i, j) + \hat{\mathbb{E}} \left[ \hat{v}_{\Pi} \left( t_{n+1}, \bar{X}_{t_{n+1}}^{t_{n}, x_{m}}, j \right) \right] \right\}$$

 $\Rightarrow$  storage of Monte Carlo sample  $\Rightarrow \mathcal{O}(N \times M)$ 

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How to	avoid the sa	ample st	orage?		

#### Memory reduction method

[Chan et al., 2004] 1-dim geometric B.M.

[Chan et al., 2006] d-dim geometric B.M.

[Chan and Wu, 2011] exponential Lévy

 $\hookrightarrow$  limited to additive processes

#### Our contribution

Extension to any Markovian process

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(Pseud	do)Random n	umber g	enerators		

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Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
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(Pseud	o)Random n	umber g	enerators		





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(Pseudo)	Random ı	number g	enerators		

set seed  $s_0 \rightarrow s_0 \rightarrow rand(M \times d)$ 



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(Pseudo	o)Random n	umber g	enerators		

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 $\rightarrow (\operatorname{rand}(M \times d)) \rightarrow s_1$ set seed  $s_0$ ► *s*<sub>0</sub>  $\left(\varepsilon_{1}^{m,k}\right)_{m=1...M}^{k=1...d}$ 

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(Pseudo	o)Random n	umber g	enerators		

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set seed 
$$s_0 \rightarrow s_0 \rightarrow rand(M \times d) \rightarrow s_1 \rightarrow rand(M \times d)$$
  
 $\left(\varepsilon_1^{m,k}\right)_{m=1...M}^{k=1...d}$ 

F		Numerical scheme	Complexity 00	Memory reduction ○●○○○○○○	Application 00000000000000	Conclusion O
	(Pseudo)	Random nı	ımber ge	nerators		

set seed 
$$s_0 \rightarrow s_0 \rightarrow rand(M \times d) \rightarrow s_1 \rightarrow rand(M \times d) \rightarrow s_2$$
  
 $\left(\varepsilon_1^{m,k}\right)_{m=1...M}^{k=1...d} \left(\varepsilon_2^{m,k}\right)_{m=1...M}^{k=1...d}$ 

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Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
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(Pseuc	lo)Random n	umber g	enerators		



• The seeds  $s_n$  can be set and stored

• The i.i.d. sequence  $\varepsilon_n^{m,\kappa}$  can be used for the Euler scheme

Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
000		00	0●000000	00000000000000	0
(Pseud	lo)Random n	umber g	enerators		

set seed 
$$s_0 \rightarrow s_0 \rightarrow rand(M \times d) \rightarrow s_1 \rightarrow rand(M \times d) \rightarrow s_2 \rightarrow \cdots$$
  
 $\left(\varepsilon_1^{m,k}\right)_{m=1...M}^{k=1...d} \left(\varepsilon_2^{m,k}\right)_{m=1...M}^{k=1...d} \cdots$ 

# The seeds s<sub>n</sub> can be set and stored The i.i.d. sequence c<sub>n</sub><sup>m,k</sup> can be used for the Euler scheme

Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
000		00	0●000000	00000000000000	0
(Pseud	o)Random n	umber g	enerators		

set seed 
$$s_0 \rightarrow s_0 \rightarrow rand(M \times d) \rightarrow s_1 \rightarrow rand(M \times d) \rightarrow s_2 \rightarrow \cdots$$
  
 $\left(\varepsilon_1^{m,k}\right)_{m=1...M}^{k=1...d} \left(\varepsilon_2^{m,k}\right)_{m=1...M}^{k=1...d} \cdots$ 

• The seeds  $s_n$  can be set and stored

• The i.i.d. sequence  $\varepsilon_n^{m,k}$  can be used for the Euler scheme

Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
000		00	00●00000	00000000000000	O
Euler sc	heme				

# Case of a diffusion

$$\begin{array}{rcl} x_{t_{n+1}} & = & F\left(x_{t_n}, \varepsilon_i\right) \\ F\left(x, \varepsilon\right) & = & x + b\left(t_n, x\right) h + \sigma\left(t_n, x\right) . \varepsilon \sqrt{h} \\ & \varepsilon_i & \sim & \mathcal{N}\left(0, 1\right) \end{array}$$

Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
000		00	00●00000	00000000000000	O
Euler sc	heme				

$$\begin{aligned} x_{t_{n+1}} &= F(x_{t_n}, \varepsilon_i) \\ F(x, \varepsilon) &= x + b(t_n, x) h + \sigma(t_n, x) . \varepsilon \sqrt{h} \\ \varepsilon_i &\sim \mathcal{N}(0, 1) \end{aligned}$$

Suppose that  $\exists F^{-1}(x,\varepsilon)$  s.t.  $\forall \varepsilon$ ,  $F(F^{-1}(x,\varepsilon),\varepsilon) = x$ 

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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion 0
General	memory re	eduction r	method (v1)		

```
Initialize X[m], m = 1..M
for i = 1..N - 1
S[i] \leftarrow \text{getseed}()
for m = 1..M
\varepsilon \leftarrow \text{rand}(d)
X[m] \leftarrow F(X[m], \varepsilon)
end
end
S[N] \leftarrow \text{getseed}()
```

#### overse Euler scheme

```
Final values in X[m], m=1..M
for i = N - 1..1
setseed (S[i])
for m = 1..M
\varepsilon \leftarrow rand(d)
X[m] \leftarrow F^{-1}(X[m], \varepsilon)
end
end
setseed (S[N])
```

Doubles the number of calls to rand ()

■ But memory complexity down to  $\mathcal{O}(M + N)$ 



```
Initialize X[m], m = 1..M
for i = 1..N - 1
S[i] \leftarrow \text{getseed ()}
for m = 1..M
\varepsilon \leftarrow \text{rand (d)}
X[m] \leftarrow F(X[m], \varepsilon)
end
end
S[N] \leftarrow \text{getseed ()}
```

#### Inverse Euler scheme

$$\begin{split} & \text{Final values in } X[m] \,, \, m = 1..M \\ & \text{for } i = N-1..1 \\ & \text{setseed} \left( S[i] \right) \\ & \text{for } m = 1..M \\ & \varepsilon \leftarrow \text{rand} \left( d \right) \\ & X[m] \leftarrow F^{-1} \left( X[m] \,, \varepsilon \right) \\ & \text{end} \\ & \text{end} \\ & \text{setseed} \left( S[N] \right) \end{split}$$

(3)

- Doubles the number of calls to rand ()
- But memory complexity down to  $\mathcal{O}(M + N)$

Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
000		00	000●0000	00000000000000	O
General	memory re	duction r	method (v1)		

```
Initialize X[m], m = 1..M
for i = 1..N - 1
S[i] \leftarrow \text{getseed}()
for m = 1..M
\varepsilon \leftarrow \text{rand}(d)
X[m] \leftarrow F(X[m], \varepsilon)
end
end
S[N] \leftarrow \text{getseed}()
```

#### Inverse Euler scheme

 $\begin{aligned} & \text{Final values in } X[m], m = 1..M \\ & \text{for } i = N - 1..1 \\ & \text{setseed} \left( S[i] \right) \\ & \text{for } m = 1..M \\ & \varepsilon \leftarrow \text{rand} \left( d \right) \\ & X[m] \leftarrow F^{-1} \left( X[m], \varepsilon \right) \\ & \text{end} \\ & \text{end} \\ & \text{setseed} \left( S[N] \right) \end{aligned}$ 

- Doubles the number of calls to rand ()
- But memory complexity down to  $\mathcal{O}(M + N)$

Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
000		00	000●0000	00000000000000	O
General	memory re	duction r	method (v1)		

```
Initialize X[m], m = 1..M
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S[i] \leftarrow \text{getseed}()
for m = 1..M
\varepsilon \leftarrow \text{rand}(d)
X[m] \leftarrow F(X[m], \varepsilon)
end
end
S[N] \leftarrow \text{getseed}()
```

#### Inverse Euler scheme

 $\begin{aligned} & \text{Final values in } X[m], m = 1..M \\ & \text{for } i = N - 1..1 \\ & \text{setseed} \left( S[i] \right) \\ & \text{for } m = 1..M \\ & \varepsilon \leftarrow \text{rand} \left( d \right) \\ & X[m] \leftarrow F^{-1} \left( X[m], \varepsilon \right) \\ & \text{end} \\ & \text{end} \\ & \text{setseed} \left( S[N] \right) \end{aligned}$ 

- Doubles the number of calls to rand ()
- But memory complexity down to  $\mathcal{O}(M + N)$

#### Euler scheme

for i = 1..N - 1 $S[i] \leftarrow \text{getseed}()$ for m = 1..M $\varepsilon \leftarrow \operatorname{rand}(d)$ end end

Final values in X[m], m=1..Msetseed (S[0])free (S)for i = N - 1..1for m = 1..M $\varepsilon \leftarrow rand (d)$  $X[m] \leftarrow F^{-1}(X[m], \varepsilon)$ end end

# No seed change nor storage

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 Formulation
 Numerical scheme
 Complexity
 Memory reduction
 Application
 Conclusion

 General memory reduction method (v2)

#### Euler scheme

```
for i = 1 N - 1
    S[i] \leftarrow \text{getseed}()
    for m = 1..M
        \varepsilon \leftarrow \mathrm{rand}(d)
    end
end
Initialize X[m], m = 1..M
for i = 1..N - 1
    setseed (S[N -i + 1])
    for m = 1..M
        \varepsilon \leftarrow \mathrm{rand}(d)
        X[m] \leftarrow F(X[m], \varepsilon)
    end
end
```

#### overse Euler scheme

```
Final values in X[m], m=1..M
setseed (S[0])
free (S)
for i = N - 1..1
for m = 1..M
\varepsilon \leftarrow \text{rand}(d)
X[m] \leftarrow F^{-1}(X[m], \varepsilon)
end
end
```

 No seed change nor storage

A (10) N (10)

#### Euler scheme

for i = 1 N - 1 $S[i] \leftarrow \text{getseed}()$ for m = 1..M $\varepsilon \leftarrow \mathrm{rand}(d)$ end end Initialize X[m], m = 1..Mfor i = 1..N - 1setseed (S[N -i + 1]) for m = 1..M $\varepsilon \leftarrow \mathrm{rand}(d)$  $X[m] \leftarrow F(X[m], \varepsilon)$ end end

#### Inverse Euler scheme

Final values in X[m], m=1..Msetseed (S[0]) free (S) for i = N - 1..1for m = 1..M $\varepsilon \leftarrow rand (d)$  $X[m] \leftarrow F^{-1}(X[m], \varepsilon)$ end end

# No seed change nor storage

A (10) < A (10) < A (10) </p>



```
for i = 1 N - 1
    S[i] \leftarrow \text{getseed}()
    for m = 1 M
        \varepsilon \leftarrow \mathrm{rand}(d)
    end
end
Initialize X[m], m = 1..M
for i = 1..N - 1
    setseed (S[N-i+1])
    for m = 1..M
        \varepsilon \leftarrow \mathrm{rand}(d)
        X[m] \leftarrow F(X[m], \varepsilon)
    end
end
```

#### Inverse Euler scheme

Final values in X[m], m=1..Msetseed (S[0]) free (S) for i = N - 1..1for m = 1..M $\varepsilon \leftarrow rand (d)$  $X[m] \leftarrow F^{-1}(X[m], \varepsilon)$ end end

 No seed change nor storage

A (10) < A (10) < A (10) </p>



$$f\left(f^{-1}\left(x
ight)
ight)=x$$
?



Formulation 000	Numerical scheme	Complexity 00	Memory reduction 000000●0	Application 00000000000000	Conclusion O
Stability	?				

$$f\left(f^{-1}\left(x\right)\right) = x + \epsilon Z, \ \epsilon > 0, \ Z \text{ r.v.}$$

Examples

Arithmetic B.M.  $\tilde{x}_{t_0} - x_{t_0} = \epsilon \sum_{i=0}^{N-1} z_i$ 

 $\hookrightarrow \mathsf{tame}$ 

Ornstein-Uhlenbeck 
$$\tilde{x}_{t_0} - x_{t_0} = \epsilon \sum_{i=0}^{N-1} \frac{1}{(1-\alpha h)^i} z_i$$
  
 $\hookrightarrow$  problem if  $T > \frac{1}{\alpha} \ln \left(\frac{1}{\epsilon}\right)$ 

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Formulation	Numerical scheme	Complexity	Memory reduction	Application	Conclusion
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Stabilizi	ng correctio	on			



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Outline					

- **1** Formulation
- 2 Numerical scheme
- 3 Complexity
- 4 Memory reduction
- 5 Application
- 6 Conclusion

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Application to i	nvestments	in power ger	neration	

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Formulatio	n Numerical scheme 00000000	Complexity 00	Memory reduction	Application •00000000000000	Conclusion O
Appl Feature	ication to inve <sup>s</sup>	stments i	n power ger	neration	
N	ot included				

- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application •0000000000000	Conclusion O
Applica Features	tion to inve	stments i	n power ger	neration	

- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application •0000000000000	Conclusion O
Applica <sup>-</sup> Features	tion to inve	stments i	in power ger	neration	

- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application •0000000000000	Conclusion 0
Applicati Features	on to invest	ments in	power gene	ration	

- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

# Included Multiple factors : Power price structural model :



Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application •0000000000000	Conclusion 0
Applicat Features	ion to inves	tments ir	n power gen	eration	

- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

# Included

- Multiple factors :
  - Demand
  - Installed capacities
  - Random outages
  - Energy prices

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application ●0000000000000	Conclusion 0
Applicat Features	ion to invest	tments in	power gene	eration	

- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

# Included

- Multiple factors :
  - Demand
  - Installed capacities
  - Random outages
  - Energy prices

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application ●0000000000000	Conclusion 0
Applicat Features	ion to invest	tments in	power gene	eration	

- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

# Included

- Multiple factors :
  - Demand
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- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

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  - Demand
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  - Energy prices

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application ●00000000000000	Conclusion O
Applicat Features	ion to inves	tments ir	n power gene	eration	

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  - Demand
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  - Energy prices

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application •0000000000000	Conclusion 0
Applicat Features	ion to inves	tments ir	n power gen	eration	

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# Included

- Multiple factors :
  - Demand
  - Installed capacities
  - Random outages
  - Energy prices

- Peaks / Scarcity
- Investments feedback

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application •0000000000000	Conclusion 0
Applicat Features	ion to inves	tments ir	n power gen	eration	

- Multiple agents / Strategic behaviour
- Multizone / Network
- Dynamic constraints / Time to build

# Included

- Multiple factors :
  - Demand
  - Installed capacities
  - Random outages
  - Energy prices

- Peaks / Scarcity
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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application ●0000000000000	Conclusion 0
Applicat Features	ion to invest	tments in	power gene	eration	

- Multiple agents / Strategic behaviour
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# Included

- Multiple factors :
  - Demand
  - Installed capacities
  - Random outages
  - Energy prices

- Peaks / Scarcity
- Investments feedback

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 000000000000000	Conclusion O
Factors	modelling				

#### Demand $D_t$

#### deterministic seasonalities + O.U. process

#### Capacities C

Controlled installed capacities :  $I_t = I_{0-} + \sum_{\tau_n \leq t} \zeta^n$ Available capacities :  $C_t = I_t \times A_t$  $A_t = \text{availability rate (outages)}$  $A_t = \mathcal{T} (\text{O.U.} + \text{seasonalities})$  $\mathcal{T} : \mathbb{R} \mapsto ]0,1[$  (cf. [Wagner, 2012])

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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Factors	modelling				

#### Demand $D_t$

deterministic seasonalities + O.U. process

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Controlled installed capacities :  $I_t = I_{0-} + \sum_{\tau_n \leq t} \zeta^n$ Available capacities :  $C_t = I_t \times A_t$  $A_t = \text{availability rate (outages)}$  $A_t = \mathcal{T} (\text{O.U.} + \text{seasonalities})$  $\mathcal{T} : \mathbb{R} \mapsto [0, 1]$  (cf. [Wagner, 2012])

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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion 0
Factors	modelling				

# Fuels and $CO_2$ prices

$$S_t^0 = CO_2$$
 price  
 $S_t^k$  =fuel price for tech.  $k \ge 1$ 

$$dS_t = \Xi S_t dt + \Sigma S_t dW_t^S$$

# $\Rightarrow$ Cointegrated geometric B.M.

- Non-stationary individual prices
- Stationary linear combinations

(provided  $rank(\Xi)$  not full, cf. [Benmenzer et al., 2007])



Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Factors	modelling				

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- Non-stationary individual prices
- Stationary linear combinations

(provided  $rank(\Xi)$  not full, cf. [Benmenzer et al., 2007])

### Full fuel price in €/MWh

$$\tilde{S}_t^k = h_k^0 S_t^0 + h_k S_t^k$$

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Power p	orice				



Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Power p	orice				



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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 0000000000000	Conclusion O
Power p	orice				



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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 0000000000000	Conclusion O
Power p	orice				



Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion 0
Objecti	ve function				

# New plant $\zeta^j$ of type j

$$\begin{array}{l} \text{Cost } \kappa^{j}\left(\zeta^{j}\right) = \kappa_{f}^{j} + \zeta^{j} \times \kappa_{p}^{j} \\ \text{Output } \mathcal{O}_{t}^{j} = \min\left(\zeta^{j} \times A_{t}^{j}, \left(D_{t} - \overline{C}_{t}^{(j-1)}\right)^{+}\right) \\ \text{Revenue } \left(P_{t} - \tilde{S}_{t}^{j}\right)^{+} - \kappa_{m}^{j} \\ \text{fotal Gain } \int_{0}^{\infty} e^{-\rho s} \left(\mathcal{O}_{s}^{j}\left(P_{s} - \widetilde{S}_{s}^{j}\right)^{+} - \kappa_{m}^{j}\right) ds - \kappa^{j} \left(\mathcal{O}_{s}^{j}\left(P_{s} - \widetilde{S}_{s}^{j}\right)^{+} - \kappa^{j} \left(P_{s} - \widetilde{S}_{s}^{j}\right) ds - \kappa^{j} \left(\mathcal{O}_{s}^{j}\left(P_{s} - \widetilde{S}_{s}^{j}\right)^{+} - \kappa^{j} \left(P_{s} - \widetilde{S}_{s}^{j}\right) ds -$$



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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion 0
Objecti	ve function				

New plant  $\zeta^{j}$  of type jCost  $\kappa^{j} \left( \zeta^{j} \right) = \kappa_{f}^{j} + \zeta^{j} \times \kappa_{p}^{j}$ Output  $O_{t}^{j} = \min \left( \zeta^{j} \times A_{t}^{j}, \left( D_{t} - \overline{C}_{t}^{(j-1)} \right)^{+} \right)$ Revenue  $\left( P_{t} - \widetilde{S}_{t}^{j} \right)^{+} - \kappa_{m}^{j}$ Total Gain  $\int_{0}^{\infty} e^{-\rho s} \left( O_{s}^{j} \left( P_{s} - \widetilde{S}_{s}^{j} \right)^{+} - \kappa_{m}^{j} \right) ds - \kappa^{j} \left( \zeta^{j} \right)$ 


Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Objecti	ve function				

New plant 
$$\zeta^{j}$$
 of type  $j$   
Cost  $\kappa^{j} (\zeta^{j}) = \kappa_{f}^{j} + \zeta^{j} \times \kappa_{p}^{j}$   
Output  $O_{t}^{j} = \min \left( \zeta^{j} \times A_{t}^{j}, \left( D_{t} - \overline{C}_{t}^{(j-1)} \right)^{+} \right)$   
Revenue  $\left( P_{t} - \tilde{S}_{t}^{j} \right)^{+} - \kappa_{m}^{j}$   
Total Gain  $\int_{0}^{\infty} e^{-\rho s} \left( O_{s}^{j} \left( P_{s} - \tilde{S}_{s}^{j} \right)^{+} - \kappa_{m}^{j} \right) ds - \kappa^{j} (\zeta^{j})$ 

$$V^{-1} = \sup_{\alpha \in \mathcal{A}_{t,i}} \mathbb{E}\left[\sum_{j=1}^{d'} \int_{t}^{\infty} e^{-\rho s} \left(O_{s}^{j} \left(P_{s} - \widetilde{S}_{s}^{j}\right)^{+} \kappa_{m}^{j}\right) ds - \sum_{\tau_{n} \geq t} e^{-\rho \tau_{n}} \kappa\left(\zeta^{j}\right)\right]$$

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Objecti	ve function				

New plant 
$$\zeta^{j}$$
 of type  $j$   
Cost  $\kappa^{j} (\zeta^{j}) = \kappa_{f}^{j} + \zeta^{j} \times \kappa_{p}^{j}$   
Output  $O_{t}^{j} = \min \left( \zeta^{j} \times A_{t}^{j}, \left( D_{t} - \overline{C}_{t}^{(j-1)} \right)^{+} \right)$   
Revenue  $\left( P_{t} - \widetilde{S}_{t}^{j} \right)^{+} - \kappa_{m}^{j}$   
Total Gain  $\int_{0}^{\infty} e^{-\rho s} \left( O_{s}^{j} \left( P_{s} - \widetilde{S}_{s}^{j} \right)^{+} - \kappa_{m}^{j} \right) ds - \kappa^{j} (\zeta^{j})$ 

$$V = \lim_{\alpha \in \mathcal{A}_{t,i}} \mathbb{E}\left[\sum_{j=1}^{d'} \int_{t}^{\infty} e^{-\rho s} \left(O_{s}^{j} \left(P_{s} - \widetilde{S}_{s}^{j}\right)^{+} \kappa_{m}^{j}\right) ds - \sum_{\tau_{n} \geq t} e^{-\rho \tau_{n}} \kappa\left(\zeta^{j}\right)\right]$$

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Objecti	ve function				

New plant 
$$\zeta^{j}$$
 of type  $j$   
Cost  $\kappa^{j} (\zeta^{j}) = \kappa_{f}^{j} + \zeta^{j} \times \kappa_{p}^{j}$   
Output  $O_{t}^{j} = \min \left( \zeta^{j} \times A_{t}^{j}, \left( D_{t} - \overline{C}_{t}^{(j-1)} \right)^{+} \right)$   
Revenue  $\left( P_{t} - \tilde{S}_{t}^{j} \right)^{+} - \kappa_{m}^{j}$   
Total Gain  $\int_{0}^{\infty} e^{-\rho s} \left( O_{s}^{j} \left( P_{s} - \tilde{S}_{s}^{j} \right)^{+} - \kappa_{m}^{j} \right) ds - \kappa^{j} (\zeta^{j})$ 

Value function  

$$v(t,x,i) = \sup_{\alpha \in \mathcal{A}_{t,i}} \mathbb{E} \left[ \sum_{j=1}^{d'} \int_{t}^{\infty} e^{-\rho s} \left( O_{s}^{j} \left( P_{s} - \widetilde{S}_{s}^{j} \right)^{+} \kappa_{m}^{j} \right) ds - \sum_{\tau_{n} \geq t} e^{-\rho \tau_{n}} \kappa(\zeta^{j}) \right]$$

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion O
Objecti	ve function				

New plant 
$$\zeta^{j}$$
 of type  $j$   
Cost  $\kappa^{j} (\zeta^{j}) = \kappa_{f}^{j} + \zeta^{j} \times \kappa_{p}^{j}$   
Output  $O_{t}^{j} = \min \left( \zeta^{j} \times A_{t}^{j}, \left( D_{t} - \overline{C}_{t}^{(j-1)} \right)^{+} \right)$   
Revenue  $\left( P_{t} - \tilde{S}_{t}^{j} \right)^{+} - \kappa_{m}^{j}$   
Total Gain  $\int_{0}^{\infty} e^{-\rho s} \left( O_{s}^{j} \left( P_{s} - \tilde{S}_{s}^{j} \right)^{+} - \kappa_{m}^{j} \right) ds - \kappa^{j} (\zeta^{j})$ 

Value function  $v(t,x,i) = \sup_{\alpha \in \mathcal{A}_{t,i}} \mathbb{E}\left[\sum_{j=1}^{d'} \int_{t}^{\infty} e^{-\rho s} \left(O_{s}^{j} \left(P_{s} - \widetilde{S}_{s}^{j}\right)^{+} \kappa_{m}^{j}\right) ds - \sum_{\tau_{n} \geq t} e^{-\rho \tau_{n}} \kappa(\zeta^{j})\right]$ 

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 000000000000000	Conclusion O
Numeric	cal example				

## 2 technologies

"Base" tech. 
$$\tilde{S}_0^1 = 40$$
€/MWh, vol.5%,  
 $I_0^1 = 67$ GW,  $\kappa_p^1 = 2.00 \, 10^9$ €/GW  
"Peak" tech.  $\tilde{S}_0^2 = 80$ €/MWh, vol.15%,  
 $I_0^1 = 33$ GW,  $\kappa_p^2 = 0.24 \, 10^9$ €/GW  
 $\tilde{S}_t^2 - 2\tilde{S}_t^1$  stationary,  $D_0 = 70$ GW

## Numerical parameters

T = 40 years (+20 years for terminal values) h = 1/730 (2 steps per day), but investments only once a year  $b = 2^6 = 64$  base functions, piecewise linear M = 5000 M.C. trajectories









Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 000000000000000	Conclusion 0
Price de	ensities				



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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 000000000000000000000000000000000000	Conclusion O
Strateg	ies comparis	son			



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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 0000000000000	Conclusion O
Extensio	ons				

- $\blacksquare More technologies \Rightarrow more dimensions$
- Asymptotic confidence intervals ⇒ second algorithm
- Multiple agents

-

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 0000000000000	Conclusion O
Extensio	ons				

# $\blacksquare \ \ \ \ More \ \ technologies \Rightarrow more \ \ dimensions$

 $\blacksquare$  Asymptotic confidence intervals  $\Rightarrow$  second algorithm

Multiple agents

-

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 0000000000000	Conclusion O
Extensio	ons				

- $\blacksquare More technologies \Rightarrow more dimensions$
- Asymptotic confidence intervals  $\Rightarrow$  second algorithm
- Multiple agents

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 0000000000000	Conclusion O
Extensio	ons				

- $\blacksquare \ \ \ \ More \ technologies \Rightarrow more \ \ dimensions$
- $\blacksquare$  Asymptotic confidence intervals  $\Rightarrow$  second algorithm

Multiple agents

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion ○
Outline					

- **1** Formulation
- 2 Numerical scheme
- 3 Complexity
- 4 Memory reduction
- **5** Application

# 6 Conclusion

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Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion •		
Conclusion							

# Numerical scheme for optimal switching

- Rate of convergence
- Memory reduction method

# Application to investments in electricity generation

- Structural power price model
- Numerical resolution

# Extensions

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion •			
Conclusion								

- Numerical scheme for optimal switching
  - Rate of convergence
  - Memory reduction method
- Application to investments in electricity generation
  - Structural power price model
  - Numerical resolution

## Extensions

Formulation 000	Numerical scheme	Complexity 00	Memory reduction	Application 00000000000000	Conclusion •			
Conclusion								

- Numerical scheme for optimal switching
  - Rate of convergence
  - Memory reduction method
- Application to investments in electricity generation
  - Structural power price model
  - Numerical resolution

# Extensions

# Questions



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